Cleanroom Software Engineering

- Harlan Mills (Linger, Dyer, Poore), IBM, 1980
- Analogy with electronic component manufacture
- Use of statistical process control features
- Certified software reliability
- Improved productivity; zero defects at delivery

Key Features

- Usage scenarios; statistical modeling
- Incremental development and release
- Separate development and acceptance testing
- No unit testing or debugging
 - Instead, formal reviews with verification conditions

Cleanroom Projects

Year	Applied technologies	Implementation	Results + No failure ever found + Programmer received gold medal from Baldridge			
1980	Stepwise refinement Functional verification	Census, 25 KLOC (Pascal)				
1983	Functional verification	Wheelwriter, 63 KLOC, three processors	 Millions of users No failure ever found 			
1980s	Functional verification Inspections	Space shuttle, 500 KLOC	 Low defect over entire function No defect in any flight Work received NASA's Quality Award 			
1987	Cleanroom engineering	Flight control, 33 KLOC (Jovial) three increments-	 Completed ahead of schedulc 2.5 errors/KLOC before any execution Error-fix effort reduced by a factor of five 			
1988	Cleanroom engineering	Commercial product, 80 KLOC (PL/I)	 Certification testing failure rate of 3.4 failures/KLOC Deployment failures of 0.1/KLOC Productivity of 740 lines/man-month 			
1989	Partial Cleanroom engineering	Satellite control, 30 KLOC (Fortran) 7	 Certification testing error rate of 3.3 failures/KLOC 50-percent improvement in quality Productivity of 780 lines/man-month 80-percent improvement in productivity 			
1990	Cleanroom engineering with reuse and new Ada design language	Research project, 12 KLOC (Ada and ADL)	 Certified to 0.9978 with 989 test cases; 36 failures found during certification (20 logic errors, or 1.7 errors/KLOC 			

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Defect Rates

- Traditional
 - Unit testing: 25 faults / KLOC
 - System testing: 25 / KLOC
 - Inspections: 20 50 / KLOC
- Cleanroom
 - < 3.5 / KLOC delivered
 - Average 2.7 / KLOC between first execution and delivery

Basic Technologies

- 1. Incremental Development
- 2. Box-Structured Specification
- 3. Function-theoretic verification
- 4. Statistical usage testing

1. Incremental Development

- Typical system < 100KLOC
- Increment: 2 15KLOC
- Team size < 14
- Each increment *End-to-End*
- Overlapped development of increments
- 12 18 weeks from beginning of specification to end of test
- Partitioning is difficult and critical

2. Formal Specification

- Box-structured design
 - Black box: stimulus-response
 - State box: formal model of system state
 - Clear box: hierarchical refinement
- Program functions
- Verification properties of control structures

Box-Structured Specification and Design

- Black Box: stimulus / condition / response; organized into tasks; Z has been used for specification; top-down, stepwise refinement; concurrency supported
- State Box: data / history view; model oriented
- **Clear Box**: procedural control (sequence, alternation, iteration, concurrent; contains nested black boxes)
- Box Definition language

State Boxes (Model-based Formal Specification)

- Description of system state in terms of *domains* (data structures without memory limitations
 - Sets, sequences, records, lists, maps, relations
- Specification of state invariant
- Specification of operations
 - Name
 - Arguments with domains
 - Validity condition (precondition)
 - Effect on state (postcondition)
- Each operation must maintain the invariant

3. Function-Theoretic Verification

 In Cleanroom, constructed programs can be checked by a parser for syntax errors, but may not be executed by the developer

– No debugging \Rightarrow cheap and predictable

- Verification is performed by a team review driven by a set of *verification conditions*
 - Questions to ask about the program code
 - Specific questions are asked about each kind of control structure
- Productivity: 3 5 x improvement in verification over debugging

Formal Inspections

- Although program proving is always an option, this involves intensive work requiring mathematical sophistication
- An alternative, used by Cleanroom software engineering, is to structure a team code inspection in terms of program functions and verification conditions and then undertake an informal review confirming all verification conditions are satisfied

Functional Verification Steps

- 1. Starting condition: program is specified by pre and post conditions
- 2. Program is parsed into prime programs
 - Prime program decomposition: parse program control flow into nested single entry/exit constructs (SESEs)
 - Usual SESEs are sequence, conditional, iteration
- 3. Proceeding top down, determine the program function for all SESEs
 - *Program function*: Description of the function of a prime program
 - Assertion placed before and after each SESE
- 4. Define verification conditions for each program point
 - Verification Conditions: things to check for each SESE
- 5. Inspect, answering all verification conditions

Program Function

- Conditions under which the program can legally execute (preconditions)
- Expression of the effect of program execution on the state of the system (postconditions)
- Expressed in terms of the program's input arguments, return value, instance variables, global variables, and side effects on the environment (disk writes, printing, etc.) but not local program variables

Program Parse

- Modern programming languages support the concept of nested blocks
 - A block is normally enclosed in braces or keyword pairs (begin-end)
- In structured programs (programs without GOTO statements), the nesting is always well formed
 - That is, there is only ever one way for control to enter the block and one way to exit. That is, they have the property of being single-entry, single exit (SESE)
 - Programs with GOTOS can be handled using special methods
- The process of determining the SESEs for a program involves parsing its control flow graph.

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Typical SESEs

Conditional Composition

Sequential Composition

Iterative Composition

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Composition of SESEs

- Each SESE can be thought of as being itself a small program with its own program function
- The overall program function is the logical composition of the program functions of its constituent SESEs
- The lowest level SESE is the single assignment statement

Verification Conditions

- If we were proving a program correct, we would construct the proof by composing the proofs of each of the SESEs
- Instead of a proof, Cleanroom uses an informal review that examines each program statements to determine its logical validity
- In particular, each type of statement has a set of questions that should be asked about it every time that it occurs in the program
- There are three ways of composing SESEs
 - Sequence, conditional and iteration

Sequence

- The simplest control structure is a sequence of two other statements or control structures
- There is one verification condition per sequence:
 - Do the constituent statements together accomplish the sequence's goal?
- This idea can readily be extended to three or more constituent statements

Sequential Composition

 Is the post assertion of the sequence equivalent to the logical composition of the first part followed by second part?

Conditional

- An if-then-else has two arms
 - Does each arm acting by itself accomplish the control structure's post condition, assuming the control structure's precondition and that the tested condition is true (or false)?
- If-then is treated as if-then-else with a null arm

Conditional

- 2.Does taking the true branch imply the post assertion?
 - The predicate of the conditional can be assumed to be true

3.Does taking the false branch imply the post assertion?

– The predicate can be assumed to be ${\tt false}$

Iteration

- There are three questions to ask about an iterative construct such as a while loop:
 - Does it terminate in all circumstances?
 - Does it accomplish its purpose when it does not execute?
 - Does it accomplish its purpose when its body is executed followed by its own execution?
- for loops and repeat loops can be defined in terms of while loops

Iteration

4. Does the loop terminate?

5.If the predicate is false, is the post assertion equivalent to the pre assertion?

6.If the predicate is true, is the post assertion of the loop equivalent to the post assertion of the body followed by the post assertion of the loop?

– Recursive!

- You may assume the predicate is true

Implications

- As teams become more experience in Cleanroom, then begin to write their programs more directly
- This typically results in very small program segments with few control structures each
- Example: 3300 lines ⇒ 600 control structures, 1000 correctness conditions

4. Statistical Usage Testing

- Certification of reliability
- Process control
- Cost-effective orientation
- Guidelines for test completion (desired reliability reached) or redesign (too many failures found)
- Stratification mechanism for dealing with critical situations
- But questions exist on how to feed back the results of testing to the development team

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Cost-Effective Testing

Table 2. Software failures for nine major IBM products, classified from rare to frequent.											
		Rare							Frequent		
Стоир		I	2	3	4	5	6	7	8		
MTTF (years)		5,000	1,580	500	158	50	15.8	5	1.58		
Percent failures in class for product	1	34.2	28.8	17.8	10.3	5.0	2.t	1.2	0.7		
	2	34 3	28.0	18.2	9.7	4.5	3.2	1.5	0.7		
	3	33.7	28.5	18.0	8.7	6.5	2.8	1.4	0.4		
	4	34.2	28.5	18.7	11.9	4.4	2.0	0.3	0.1		
	5	34.2	28.5	18.4	9.4	4.4	2.9	1.4	0.7		
· •	6	32.0	28.2	20.1	11.5	5.0	21	0.8	0.3		
	7	34.0	28.5	18.5	9.9	4.5	2.7	1.4	0.6		
	8	31.9	27.1	18.4	11.1	6.5	2.7	14	1.1		
	9	31.2	27.6	20.4	12.8	5.6	1.9	0.5	0.0		
Average percentage Iailures		33.4	28.2 · · · ·	18.7 	10.6	\$.2 2 +	2.5	1.0	04		
Probability of a failu for this frequency	re	0.008	() 0021	.044	0.079	0.123	0.187	0 237	0.300		

Testing Process

- Usage distribution models
 - From competitors, earlier versions, analysis
- Markov usage chain
 - State transition probability matrix
- Statistics
 - Π (proportion of time spent in each state)
 - n (number of states visited before a given state is reached)
 - s (number of tests needed to reach a state).
- Random test generation
 - Design required
- Test execution and test chain generation, including failure states
- Statistics
 - R (reliability)
 - MTBF (mean time between failures)
 - D (divergence of test chain from usage chain)

Testing Process Overview

- Usage distribution models; other software, earlier versions, analysis
- Construct Markov usage chain / probability matrix
- Computations of Π (proportion of time spent in each state), n (number of states visited before a given state is reached), and s (number of tests needed to reach a state).
- Random test generation (some design required here to deal with constraints)
- Test execution and test chain generation, including failure states
- Calculations of R (reliability), MTBF (mean time between failures), and D (divergence of test chain from usage chain) 6/18/2007 © 2007, Spencer Rugaber 28

Testing Example

- COBOL / SF parser generator
- Four increments; 120 random tests
- Last 115 executions correct
- 12 failures in first five executions
- 3.9 faults / KLOC
- No new failures in four years of use

Usage Model For Unix Mail



Results Of Independent Empirical Evaluation

- 15 3-person teams; 10 of them used Cleanroom
- 6/10 delivered 91% of functionality
- Requirements better met and less failures
- More comments, less dense control flow
- Better adherence to schedule
- Developers expressed satisfaction with process

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Results

- Defects: 2 5 / KLOC versus 10-30 / KLOC for debugging
- Productivity: 3 5 × improvement in verification over debugging
- Reliability: statistical usage testing 20 × as effective as coverage testing

Cleanroom Tools

- Test case generator
- Reliability analysis package
 - Spreadsheet
- Verification-based inspection syntax analyzer
 - Script for inspection
- Management assistant
 - Reports on process